

[S]=T. Sauer, Numerical Analysis, Second International Edition, Pearson, 2014

## "Teorioppgaver"

**1** Oppgave 10.1.1

### Solution:

(a) x = [0, 1, 0, -1],

$$y_k = \frac{1}{2} \sum_{j=0}^{3} x_j \exp(-i2\pi jk/4) = \frac{1}{2} [\exp(-i2\pi k/4) - \exp(-i6\pi k/4)]$$

Thus  $y_0 = 0, y_1 = [\exp(-i\pi/2) - \exp(-i3\pi/2)]/2 = [-i-i]/2 = -i, y_2 = [\exp(-i\pi) - \exp(-i3\pi)]/2 = [-1 - (-1)]/2 = 0, y_3 = [\exp(-i3\pi/2) - \exp(-i9\pi/2)]/2 = [i - (-i)]/2 = i,$ (b)  $x = [1, 1, 1, 1], y_0 = 2, y_1 = [\exp(-i\pi/2) + \exp(-i\pi) + \exp(-3i\pi/2) + \exp(-2i\pi)]/2 = [-i - 1 + i + 1]/2 = 0, y_2 = 0, y_3 = 0.$ (c) x = [0, -1, 0, 1], which is just negative of x in (a). Since DFT is linear  $\implies$  y = [0, i, 0, -i].(d) x = [0, 1, 0, -1, 0, 1, 0, -1], that is, the repetition of x in (a) twice. Thus  $y_k = \frac{1}{\sqrt{8}} \sum_{j=0}^7 x_j \exp(-i2\pi jk/8)$  $= \frac{1}{2\sqrt{2}} \sum_{j=0}^3 x_j \exp(-i2\pi jk/8) + \frac{1}{2\sqrt{2}} \sum_{j=4}^7 x_j \exp(-i2\pi jk/8)$ 

$$= \frac{1}{2\sqrt{2}} \sum_{j=0}^{3} x_j \exp(-i2\pi jk/8) + \frac{1}{2\sqrt{2}} \exp(-i2\pi 4k/8) \sum_{j=0}^{3} x_j \exp(-i2\pi jk/8)$$
$$= \frac{1}{2\sqrt{2}} [1 + \exp(-i\pi k)] \sum_{j=0}^{3} x_j \exp(-i2\pi jk/8).$$

In particular, the factor in the square brackets equals 2 for even k and 0 for odd k, and therefore  $y_k = 0$  for all odd k = 1, 3, 5, 7. Furthermore, for even k = 0, 2, 4, 6 the sum in the last expression equals that of the sum in (a) for k/2. Thus  $y_0 = 0$ ,  $y_2 = -i\sqrt{2}$ ,  $y_4 = 0$ ,  $y_6 = i\sqrt{2}$ .

## 2 Oppgave 10.1.7

 $\begin{array}{|c|c|c|c|c|c|c|c|} \hline \textbf{3} & \text{Prove the following identity (Parseval's theorem for DFT): Let } y^{(1)} = \text{DFT}(x^{(1)} \text{ and } \\ y^{(2)} = \text{DFT}(x^{(2)}). & \text{Then } \sum_{j=0}^{n-1} x_j^{(1)} \bar{x}_j^{(2)} = \sum_{j=0}^{n-1} y_k^{(1)} \bar{y}_k^{(2)}. \\ & \text{Hint: } y^{(1,2)} = F_n x^{(1,2)}. \end{array}$ 

#### Solution:

$$\sum_{j=0}^{n-1} y_k^{(1)} \bar{y}_k^{(2)} = [\bar{y}^{(2)}]^T y^{(1)} = [\bar{F}_n \bar{x}^{(2)}]^T F_n x^{(1)} = [\bar{x}^{(2)}]^T \underbrace{\bar{F}_n^T F_n}_{=I} x^{(1)} = \sum_{j=0}^{n-1} x_j^{(1)} \bar{x}_j^{(2)},$$

since  $\bar{F}_n^T = F_n^{-1}$  (see for example (10.11) in [S]).

4 Prove the following identity (circular shift theorem for DFT): Let  $\tilde{x}_j = \exp(i2\pi jm/n)$ , where  $i^2 = -1$  and m is an integer. Let y = DFT(x), and  $\tilde{y} = \text{DFT}(\tilde{x})$ . Show that  $\tilde{y}_k = y_{k-m}$ , where the subscript k - m is understood modulo n (in other words, the sequence y is assumed to be periodically repeating with period n).

#### Solution:

$$\tilde{y}_k = \frac{1}{\sqrt{n}} \sum_{j=0}^{n-1} \tilde{x}_j \exp(-i2\pi jk/n) = \frac{1}{\sqrt{n}} \sum_{j=0}^{n-1} x_j \exp(i2\pi jm/n) \exp(-i2\pi jk/n)$$
$$= \frac{1}{\sqrt{n}} \sum_{j=0}^{n-1} x_j \exp(-i2\pi j(k-m)/n) = y_{k-m},$$

since  $k \mapsto \exp(-i2\pi jk/n) = \cos(2\pi jk/n) - i\sin(2\pi jk/n)$  is a periodic function with period n.

# "Computeroppgaver"

**5** Implement Cooley–Tukey's algorithm for computing DFT:

```
function [y,nop] = myfft(x)
% Implementation of FFT
%
% Input: x, supposed to be a vector of length 2^p
%
% Output: y: FFT(x)
% nop: number of operations required
%
```

Verify it against Matlab's FFT (recall that Matlab uses a different scaling of DFT). Check that the number of operations needed by Cooley–Tukey's algorithm scales as  $O(N \log(N))$  by plotting the number of iterations vs. N and  $N \log(N)$  on a log-log plot for a range of  $N = 2^p$ .

## Solution:

See myfft.m and check\_myfft.m.