

# Error Resilience in the Number-Theoretic Transform

PQC acceleration for safety critical applications

Mohamed Abdelmonem

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Simula  
UiB

# Motivation

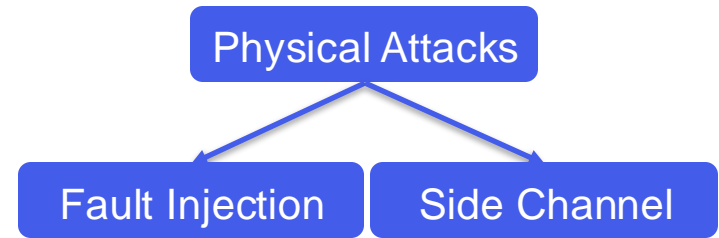


Post-Quantum Cryptography:  
Algorithms based on **classical**  
operations that resist attacks  
from quantum computers



Hard problems:

- Lattice based
- Code based
- Hash based
- Multivariate based
- Isogeny based



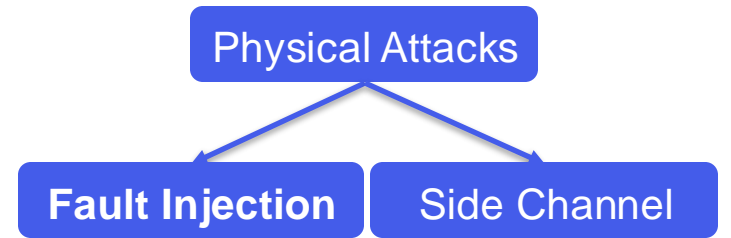
# Motivation



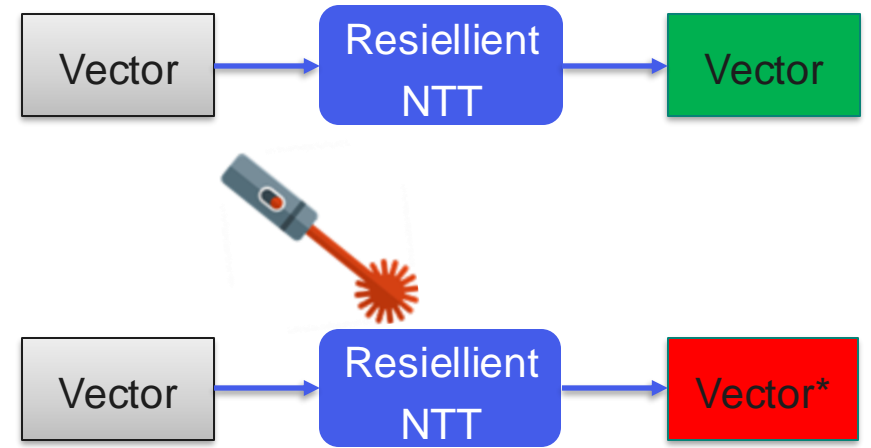
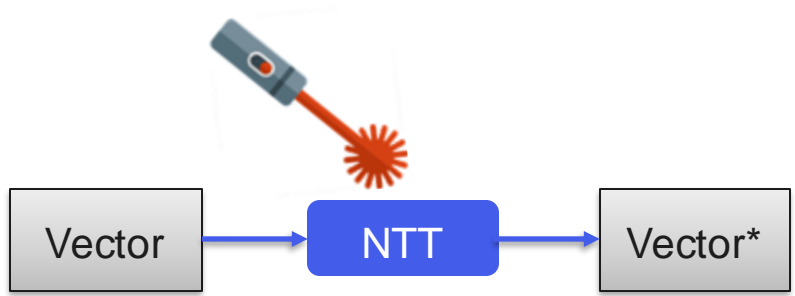
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# Example



# Number Theoretic Transform

## – Module-Lattice-Based

- Use polynomial ring  $R_q = \mathbb{Z}_q[X]/\phi(x)$  instead of  $\mathbb{Z}_q$
- Better efficiency
- Same security



Fast arithmetic on  $R_q = \mathbb{Z}_q[X]/\phi(x)$  necessary

- Schoolbook-Multiplication”  $O(n^2)$  operations
- Fast Multiplication  $O(n \log(n))$  operations

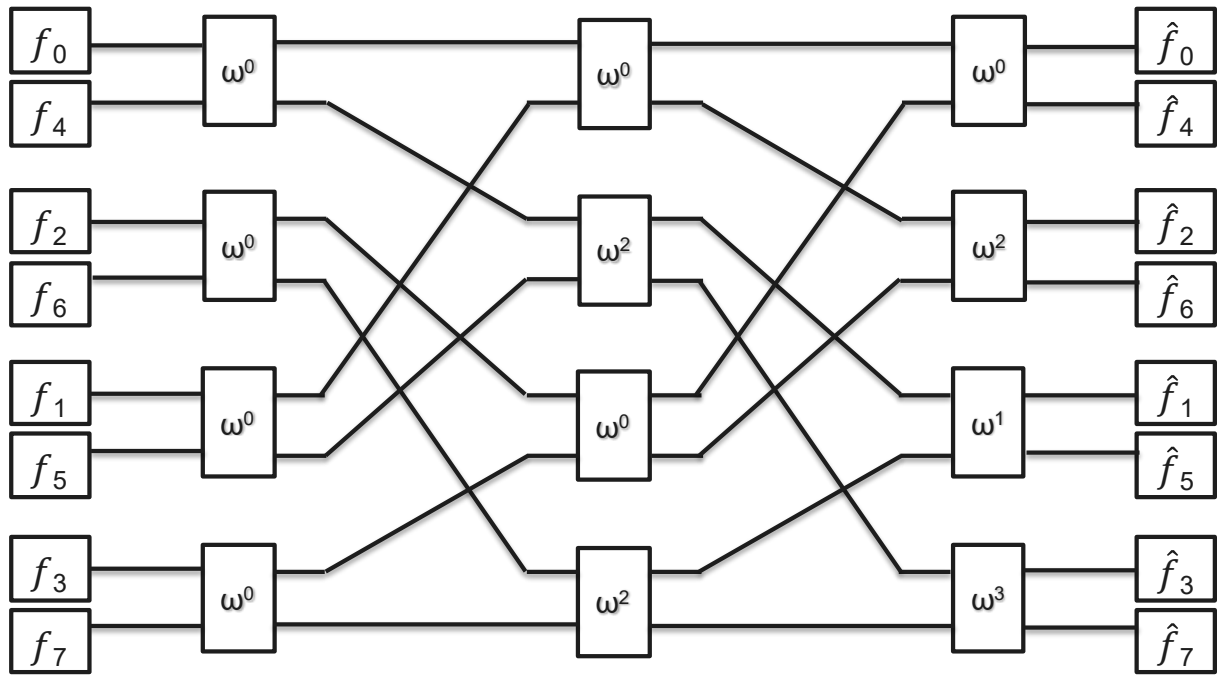
## – Fast Multiplication

Calculate product as

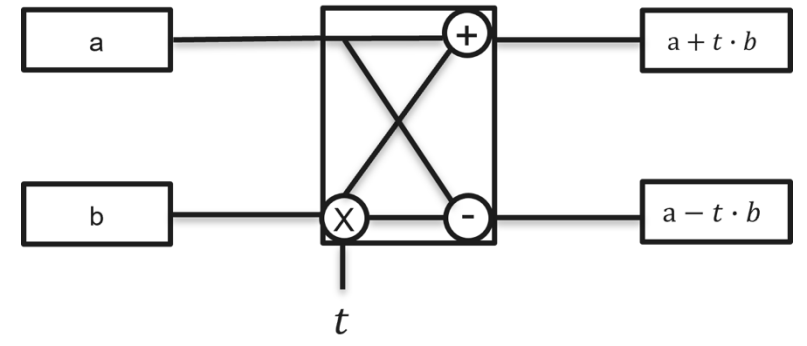
$$f \times g = NTT^{-1}(NTT(f) \circ NTT(g))$$

where  $\circ$  is the elementwise multiplication

# Number Theoretic Transform

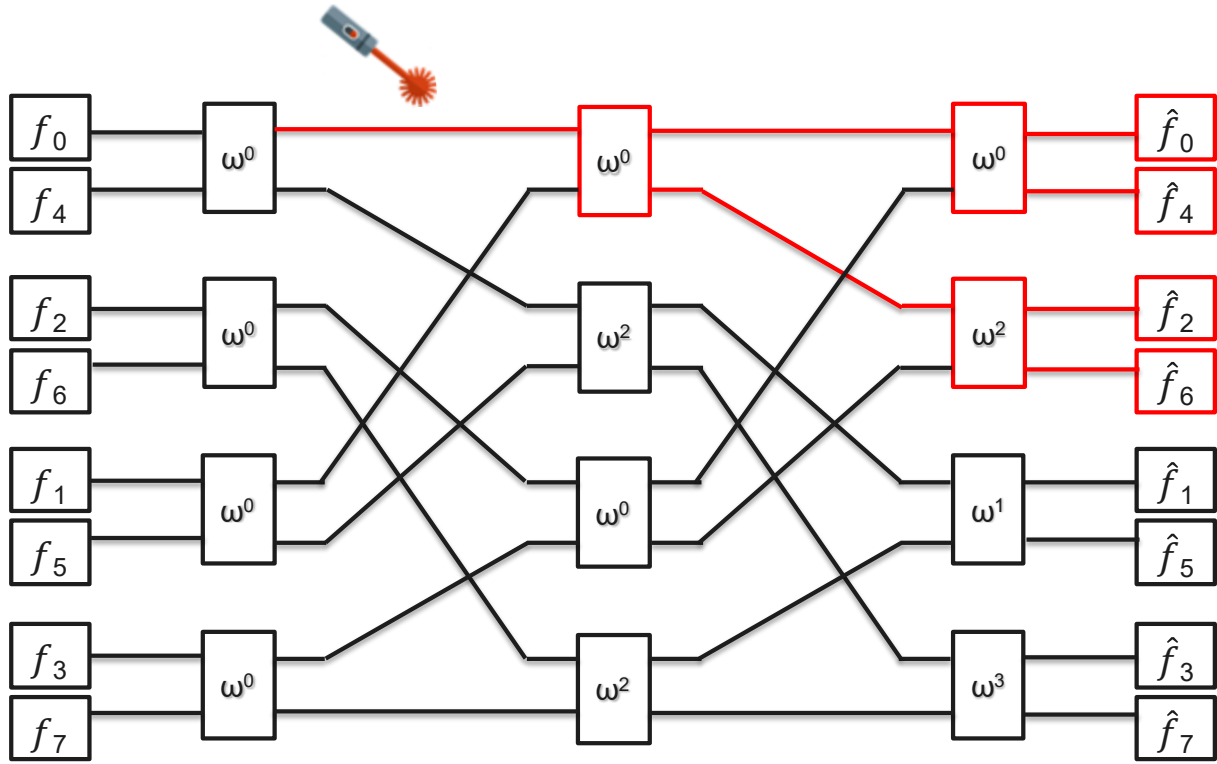


General structure of an NTT

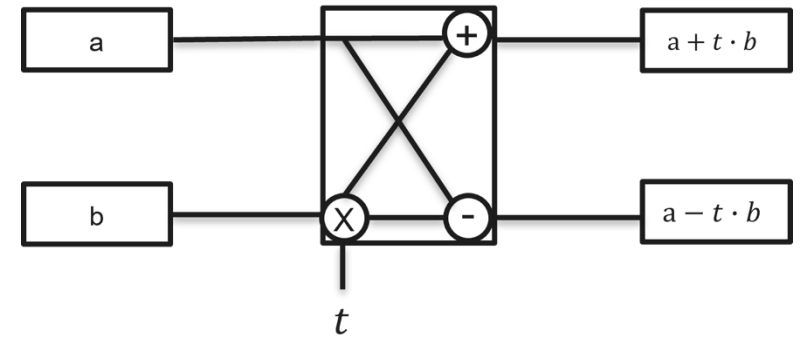


Breakdown into core operations

# Number Theoretic Transform



General structure of an NTT

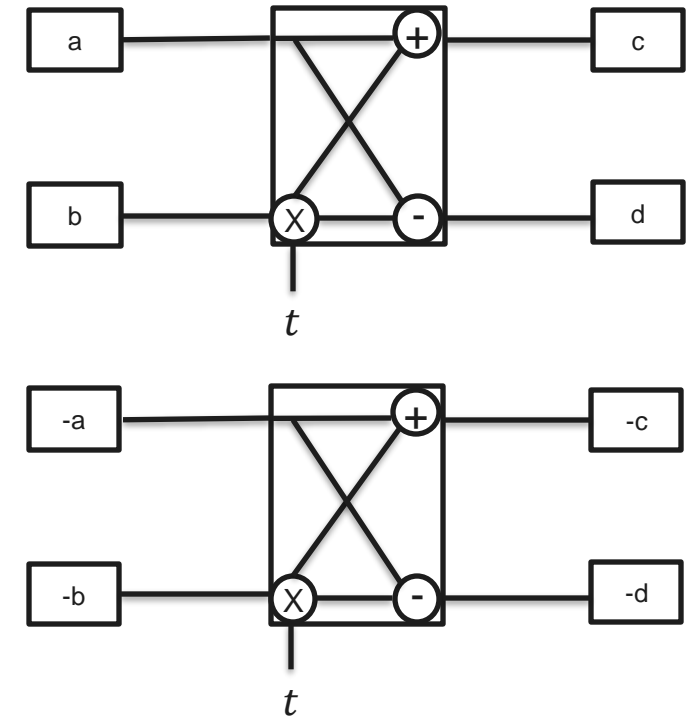
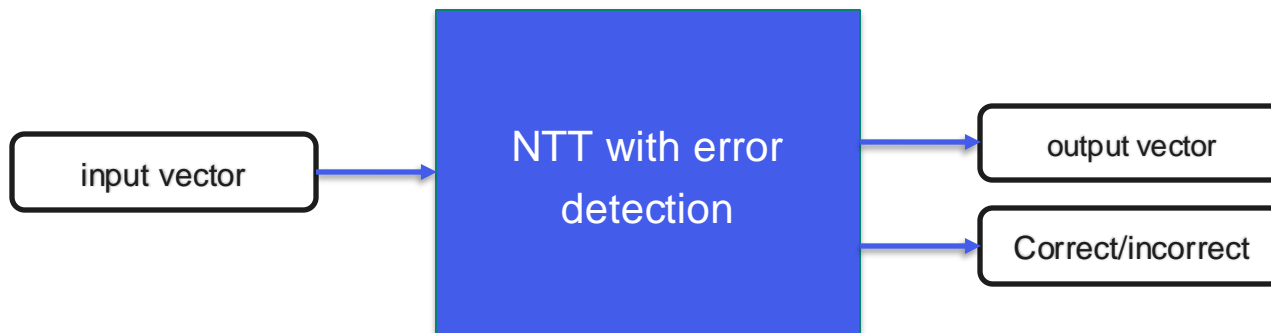


Breakdown into core operations

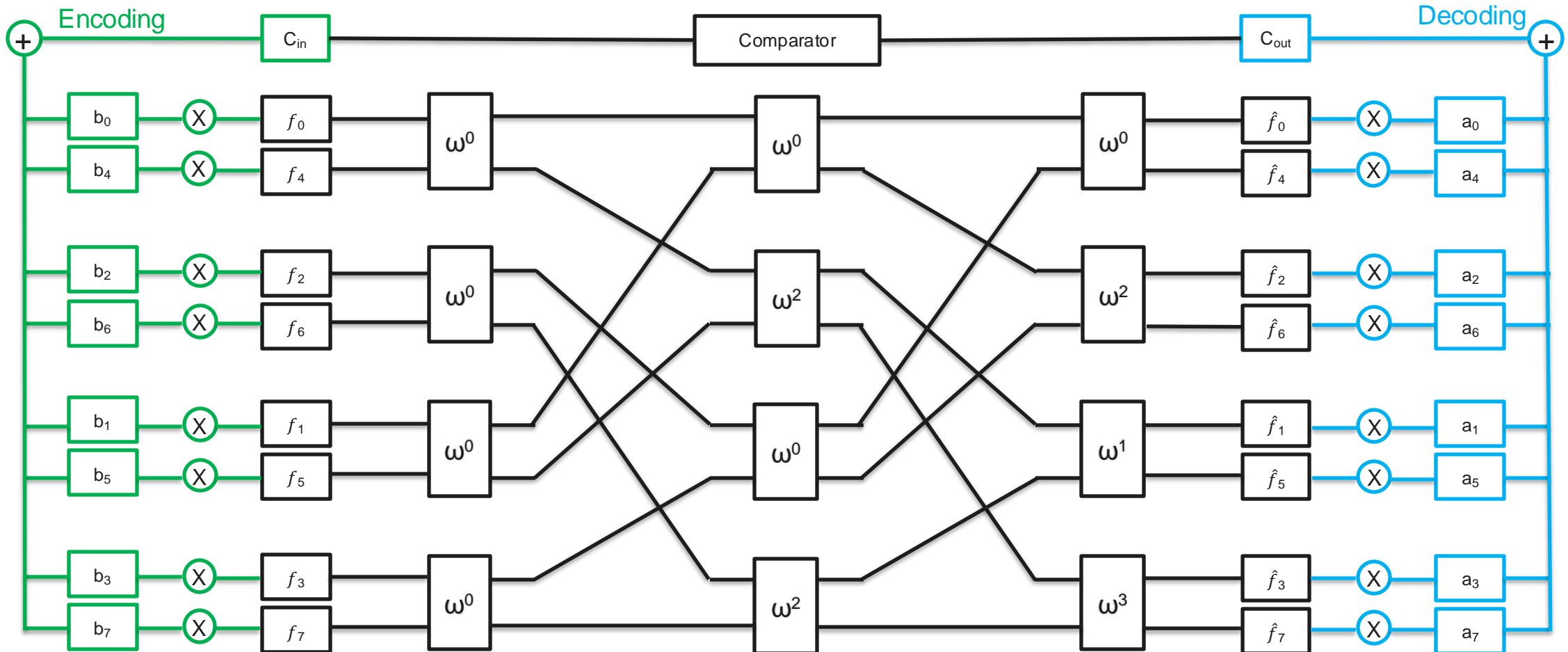
Recomputing (Sarker 2019)

## Techniques in the Literature

- Each butterfly operations is done twice → compare results
- Disadvantages:
  - high overhead
  - tests only butterfly calculations
- Our goal:
  - Test the whole NTT block with less overhead

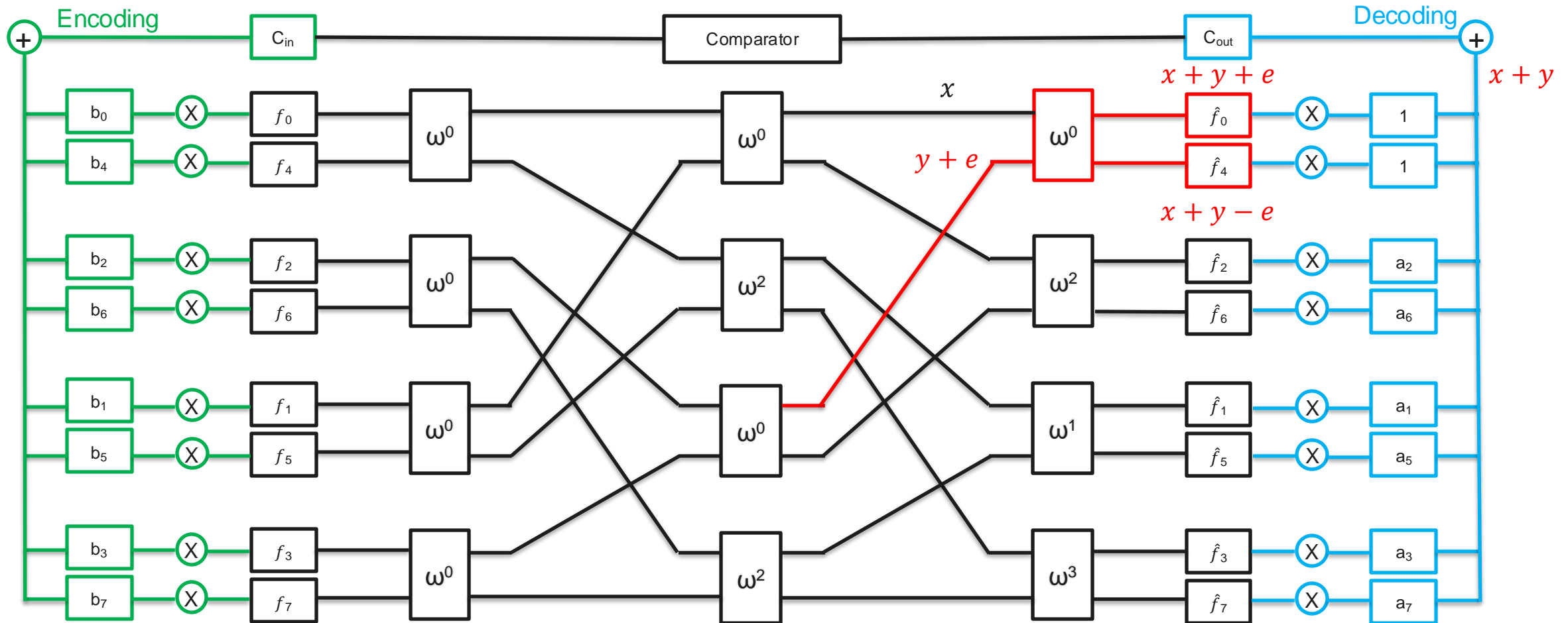


# Applying methods known for FFT





# Applying methods known for FFT



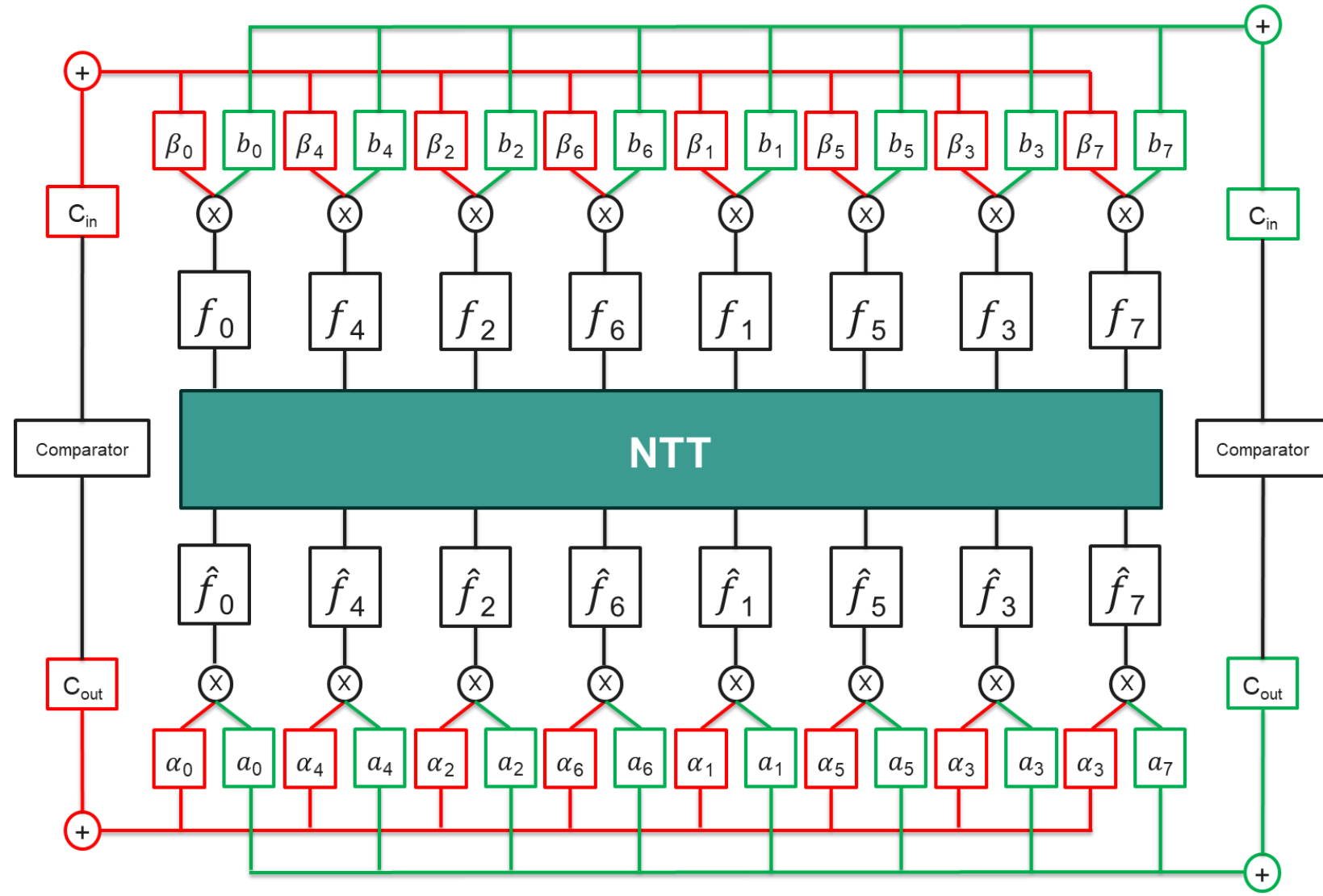
➡  $a_0$  and  $a_4$  can not be the same

## Our work

- Derive conditions on **a**'s to detect **one and two Faults in NTT network**
- One Fault:
  - show that **a** exists that meets these conditions for the parameter sets of **Dilithium, Falcon and Kyber**
  - we provide an efficient implementation that requires only  $N$  **extra multiplications** (recomputing  $N/2 \log N$  extra multiplications)

Scheme	Dilithium	Kyber	Falcon I	Falcon V
<b>Overhead</b>	0.25	0.286	0.222	0.2

# Two error detection



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Scheme	Dilithium	Kyber	Falcon I	Falcon V
<b>Overhead</b>	0.25	0.286	0.222	0.2

- Two Faults:
  - show that **a**'s exist that meet these conditions for the parameter sets of **Dilithium**
  - upper bound: about  $N^2$  restrictions on  $\mathbf{a}_i \in \mathbb{Z}_q^N$

Scheme	Ring degree $N$	$N^2$	Modulus $q$
Dilithium	256	65536	$< 8380417$
Kyber	256	65536	$> 3329$
FALCON I	512	262144	$> 12289$
FALCON V	1024	1048576	$> 12289$

## Future work

- Find **a's** that also work for Kyber and Falcon
- Error correction
- Find Fault injection attacks and countermeasures on other PQC schemes
  - 40 new NIST **Additional Digital Signature Schemes** in Round 1